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A BINARY QUANTIFIER FOR DEFINITE DESCRIPTIONS IN INTUITIONIST NEGATIVE FREE LOGIC: NATURAL DEDUCTION AND NORMALISATION

Abstract

This paper presents a way of formalising definite descriptions with a binary quantifier ι , where $\iota x[F, G]$ is read as ‘The F is G ’. Introduction and elimination rules for ι in a system of intuitionist negative free logic are formulated. Procedures for removing maximal formulas of the form $\iota x[F, G]$ are given, and it is shown that deductions in the system can be brought into normal form.

Keywords: definite descriptions, negative intuitionist free logic, natural deduction, normalization.

1. Introduction

The definite description operator ι , the formal analogue of the definite article ‘the’, is usually taken to be a term forming operator: if A is a predicate, then ιxA is a term denoting the sole A , if there is one, or nothing or an arbitrary object if there is no or more than one A . This paper follows a different approach to definite descriptions by formalising them instead with a primitive binary quantifier: ι forms a formula from two predicates, and ‘The F is G ’ is formalised as $\iota x[Fx, Gx]$. The notation, and the way of treating definite descriptions that comes with it, was suggested by Dummett [2, p.162].¹

¹Bostock considers a similar approach and explains definite descriptions as a special case of restricted quantification, where the restriction is to a single object. [1, Sec. 8.4] Bostock writes $(\iota x: Fx) Gx$ for ‘The F is G ’, but prefers to treat definite descriptions

The current paper treats definite descriptions purely proof theoretically. The proof theory of a term forming ι operator has been investigated in the context of sequent calculi for classical free logic by Indrzejczak [3, 4]. Tennant gives rules for such an operator in natural deduction [7, p.110].² The approach followed here may be new to the literature.

In this paper, I investigate the binary quantifier ι in the context of a system of natural deduction for intuitionist negative free logic. The application of the present treatment of definite descriptions to other systems of logic and their comparisons to systems known from the literature are left for further papers. To anticipate, using a negative free logic, the approach proposed here lends itself to a natural formalisation of a Russellian theory of definite descriptions, while it provides a natural formalisation of Lambert's minimal theory of definite descriptions when the logic is positive and free.

First, notation. I will use A_t^x to denote the result of replacing all free occurrences of the variable x in the formula A by the term t or the result of substituting t for the free variable x in A . t is free for x in A means that no (free) occurrences of a variable in t become bound by a quantifier in A after substitution. In using the notation A_t^x I assume that t is free for x in A or that the bound variables of A have been renamed to allow for substitution without 'clashes' of variables, but for clarity I also often mention the condition that t is free for x in A explicitly. I also use the notation Ax to indicate that x is free in A , and At for the result of substituting t for x in A .

2. Natural Deduction for ι in Intuitionist Logic

The introduction and elimination rules for the propositional logical constants of intuitionist logic **I** are:

$$\wedge I: \frac{A \quad B}{A \wedge B} \qquad \wedge E: \frac{A \wedge B}{A} \qquad \frac{A \wedge B}{B}$$

with a term forming operator. I owe the reference to Bostock to a referee for this journal, who also pointed me to the paper by Scott to be referred to in footnote 4 and made valuable comments on this paper.

²Tennant is not explicit whether the logic in this paper is classical or intuitionist. However, as he is partial to anti-realism and constructive mathematics, we are justified in assuming that his preferred route is to add these rules to a system of intuitionist free logic. The rules are also in [8, Ch. 7.10], where the logic is classical.

$$\begin{array}{l}
 \rightarrow I: \frac{\overline{A}^i \quad \Pi \quad B}{A \rightarrow B}^i \qquad \rightarrow E: \frac{A \rightarrow B \quad A}{B} \\
 \vee I: \frac{A}{A \vee B} \qquad \frac{B}{A \vee B} \qquad \vee E: \frac{\overline{A}^i \quad \Pi \quad \overline{B}^i}{A \vee B} \quad \frac{C}{C}^i \\
 \perp E: \frac{\perp}{B}
 \end{array}$$

where the conclusion B of $\perp E$ is restricted to atomic formulas.

The introduction and elimination rules for the quantifiers of **I** are:

$$\forall I: \frac{A_y^x}{\forall x A} \qquad \forall E: \frac{\forall x A}{A_t^x}$$

where in $\forall I$, y is not free in any undischarged assumptions that A_y^x depends on, and either y is the same as x or y is not free in A ; and in $\forall E$, t is free for x in A .

$$\exists I: \frac{A_t^x}{\exists x A} \qquad \exists E: \frac{\overline{A}_y^x \quad \Pi \quad C}{\exists x A}^i$$

where in $\exists I$, t is free for x in A ; and in $\exists E$, y is not free in C nor any undischarged assumptions it depends on in Π except A_y^x , and either y is the same as x or it is not free in A .

The introduction and elimination rules for identity are:

$$= I: \frac{}{t = t} \qquad = E: \frac{t_1 = t_2 \quad A_{t_1}^x}{A_{t_2}^x}$$

where A is an atomic formula. To exclude vacuous applications of $= E$, we can require that x is free in A and that t_1 and t_2 are different. An induction over the complexity of formulas shows that the rule holds for formulas of any complexity.

To formalise definite descriptions, one could add the binary quantifier ι to **I**. Its introduction and elimination rules would be:

$$\iota I : \frac{\frac{F_t^x \quad G_t^x}{\iota x[F, G]} \quad \frac{\Pi}{z = t} \quad \overline{F_z^x}^i}{\iota x[F, G]}^i$$

where t is free for x in F and in G , and z is different from x , not free in t and does not occur free in any undischarged assumptions in Π except F_z^x .³

$$\iota E^1 : \frac{\frac{\iota x[F, G]}{C} \quad \overbrace{F_z^x, G_z^x}^i}{C}^i}{C}^i$$

where z is not free in C nor any undischarged assumptions it depends on except F_z^x and G_z^x , and either z is the same as x or it is not free in F nor in G .

$$\iota E^2 : \frac{\iota x[F, G] \quad F_{t_1}^x \quad F_{t_2}^x}{t_1 = t_2}$$

where t_1 and t_2 are free for x in F .

For simplicity we could require that x occurs free in F and G . If we don't, the truth or falsity of $\iota x[F, G]$ may depend on properties of the domain of quantification: if F is true and does not contain x free, then $\iota x[F, G]$ is false if there is more than one thing in the domain of quantification, and it is true if there is only one thing and G is true (of the one thing, if x is free in G).

$\iota x[F, G]$ and $\exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)$ are interderivable. Notice that the rules for identity are not applied in the two deductions to follow.

³A more precise and general statement of the introduction rule for ι would result if we were to require Π to be a deduction of $(y = t)_z^y$ from $(F_y^x)_z^y$, where y is different from x and not free in t , and either z is the same as y or z is not free in F_y^x nor in $y = t$.

$$1. \iota x[F, G] \vdash \exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)$$

Let y be different from x and not free in F or G :

$$\frac{\frac{\frac{\frac{\frac{\iota x[F, G]}{F} \quad \overline{F_y^x}^1 \quad \overline{F}^2}{y = x} \quad \iota E^2}{(F_y^x \rightarrow y = x)}^1}{\forall y(F_y^x \rightarrow y = x)} \quad \overline{F}^2}{F \wedge \forall y(F_y^x \rightarrow y = x)} \quad \overline{G}^2}{(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)} \quad \exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)} \quad \iota E^1}{\exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)} \quad \iota E^1$$

$$2. \exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G) \vdash \iota x[F, G]$$

Let y be different from x and not free in F or G , and let \otimes be the formula $(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)$:

$$\frac{\frac{\frac{\frac{\frac{\otimes}{F} \quad \frac{\otimes}{G}}{\exists x \otimes} \quad \frac{\frac{\frac{\overline{\otimes}^2}{\forall y(F_y^x \rightarrow y = x)}}{F_y^x \rightarrow y = x} \quad \overline{F_y^x}^1}{y = x} \quad \iota I}{\iota x[F, G]} \quad \iota I}{\exists x \otimes} \quad \iota I}{\iota x[F, G]} \quad \iota I$$

3. Intuitionist Free Logic

It is more interesting to add the ι quantifier to a free logic. I will use formalisations of intuitionist free logic with a primitive predicate $\exists!$, to be interpreted as ‘ x exists’ or ‘ x refers’ or ‘ x denotes’. The introduction and elimination rules for the quantifiers are:

$$\forall I : \frac{\frac{\overline{\exists!y}^i}{\Pi} \quad A_y^x}{\forall x A}^i \quad \forall E : \frac{\forall x A \quad \exists!t}{A_t^x}$$

where in $\forall I$, y does not occur free in any undischarged assumptions of Π except $\exists!y$, and either y is the same as x or y is not free in A ; and in $\forall E$, t is free for x in A .

$$\exists I : \frac{A_t^x \quad \exists!t}{\exists x A} \quad \exists E : \frac{\exists x A \quad \frac{\overline{A_y^x}^i, \overline{\exists!y}^i}{\Pi} \quad C}{C}^i$$

where in $\exists I$, t is free for x in A ; and in $\exists E$, y is not free in C nor any undischarged assumptions of Π , except A_y^x and $\exists!y$, and either y is the same as x or it is not free in A .

The elimination rule for identity in intuitionist free logic is the same as in **I**.

In *intuitionist positive free logic* **IPF**, identity has the same introduction rule as in intuitionist logic, i.e. $\vdash t = t$, for any term t . Semantically speaking, in positive free logic any statement of self-identity is true, irrespective of whether a term refers or not.

In *intuitionist negative free logic* **INF** the introduction rule for identity is weakened and requires an existential premise:

$$= I^n : \frac{\exists!t}{t = t}$$

In **INF** the existence of t_i may be inferred if t_i occurs in an atomic formula:

$$AD : \frac{A t_1 \dots t_n}{\exists!t_i}$$

where A is an n -place predicate letter (including identity) and $1 \leq i \leq n$. Speaking semantically, for an atomic sentence, including identities, to be true, all terms in it must refer. If the language has function symbols, there is also the rule of functional denotation:

$$FD : \frac{\exists! f t_1 \dots t_n}{\exists! t_i}$$

where f is an n -place function letter and $1 \leq i \leq n$. Speaking semantically, for the value of a function to exist, all of its arguments must exist. $= I^n$, AD and FD are called the rules of strictness.⁴

Hintikka’s Law $\exists! t \leftrightarrow \exists x x = t$, where x not in t , is provable in **INF** and **IPF**. In **IPF**, it suffices to observe the following:

$$\frac{\overline{t = t}}{\exists x x = t} \quad \exists! t \qquad \frac{\overline{x = t}^1 \quad \overline{\exists! x}^1}{\exists! t}^1}{\exists x x = t} \quad \exists! t$$

In **INF**, conclude $t = t$ from $\exists! t$.

The degree of a formula is the number of connectives occurring in it. \perp , being a connective, is of degree 1. This excludes the superfluous case in which \perp is inferred from \perp by $\perp E$. $\exists! t$ is an atomic formula of degree 0.

The *major premise* of an elimination rule is the premise with the connective that the rule governs. The other premises are minor premises. A *maximal formula* is one that is the conclusion of an introduction rule and the major premise of an elimination rule for its main connective. A *segment* is a sequence of formulas of the same shape, all minor premises and conclusions of $\vee E$ or $\exists E$, except the first and the last one; the first is only a minor premise, the last only a conclusion. A segment is maximal if its first formula has been derived by an application of an introduction rule for its main connective, and its last formula is the major premise of an elimination rule. A deduction is *in normal form* if it contains neither maximal formulas nor maximal segments. A normalisation theorem establishes that any deduction can be brought into normal form by applying *reduction procedures* for the removal of maximal formulas from deductions and *permutative reduction procedures* for reducing maximal segments to maximal formulas.

Notice that the conditions imposed on applications of $= E$ have the consequence that there are no maximal formulas of the form $t_1 = t_2$.

⁴**INF** is the system introduced by Scott [6] and called **Nie** by Troelstra and Schwichtenberg [9, 200] but with a simpler theory of identity. It is the system that results if classical *reductio ad absurdum*, the rule that licenses the derivation of A if $\neg A$ entails a contradiction, is not taken to form part of the system Tennant presents in [8, Ch. 7.10].

AD and FD have the characteristics of introduction rules for $\exists!$, and $= I^m$ has the characteristics of an elimination rule for it. In a sense $\forall E$ and $\exists I$ of free logic also eliminate formulas of the form $\exists!t$. I will, however, not count these rules as introduction and elimination rules for $\exists!$, as there is no general way of removing formulas of the form $\exists!t$ that have been concluded by AD or FD and are premises of $= I^m$, $\forall E$ or $\exists I$.

Proofs of the normalisation theorem for intuitionist logic, such as those given by Prawitz [5, Ch. IV.1] and Troelstra and Schwichtenberg [9, Ch. 6.1], can be modified to carry over to the intuitionist free logics considered here.

A normalisation theorem for intuitionist negative free logic with a term forming ι operator can be reconstructed from material Tennant provides in [8]. In particular, as in the case of **I**, we can assume that every application of $\forall I$ and $\exists E$ has its own variable, that is, the free variable y of an application of such a rule occurs only in the hypotheses discharged by the rule and formulas concluded from them and, for $\forall I$, in the premise of that rule and the formulas it has been derived from. This way we avoid ‘clashes’ between the restrictions on the variables of different application of these rules when reduction procedures are applied to a deduction containing maximal formulas. Applying the reduction procedures for quantifiers of free logic can only introduce maximal formulas of lower degree than the one removed. I leave the details to the reader.

4. Natural Deduction for ι in INF

The interderivability of $\iota x[F, G]$ and $\exists x(F \wedge \forall y(F_y^x \rightarrow x = y) \wedge G)$ is the hall mark of a Russellian theory of definite descriptions, in which any statement of the form ‘The F is G ’ is false if there is no F or if there is more than one. It is the generally accepted treatment of definite descriptions in negative free logic. To establish how to modify the rules for ι given in Section 2 to yield a Russellian theory of definite descriptions when the logic is intuitionist negative free logic, we analyse the deductions establishing the interderivability of $\iota x[F, G]$ and $\exists x(F \wedge \forall y(F_y^x \rightarrow x = y) \wedge G)$ in **I** given at the end of that section.

Looking at the derivation of $\exists x(F \wedge \forall y(F_y^x \rightarrow x = y) \wedge G)$ from $\iota x[F, G]$, had the application of the universal quantifier introduction rule be one of free logic, it would have allowed the discharge of an assumption $\exists!y$, and had the existential quantifier introduction rule been one of free logic, a

further assumption $\exists!x$ would have been required. Both lend themselves as additional premises of ιE^2 , as premises analogous to the existence assumptions in the rules of the quantifiers of free logic. $\exists!y$ would be discharged by the application of the universal quantifier introduction rule of free logic, so in order for the conclusion of the deduction not to depend on $\exists!x$, it would have to be discharged, and the only option here is that it is discharged by the application of ιE^1 . This is also a natural option, corresponding, as it does, to the discharge of existence assumptions by the quantifier rules of free logic.

Generalising the first observation, we add the premises $\exists!t_1$ and $\exists!t_2$ to ιE^2 :

$$\iota E^2 : \frac{\iota x[F, G] \quad \exists!t_1 \quad \exists!t_2 \quad F_{t_1}^x \quad F_{t_2}^x}{t_1 = t_2}$$

where t_1 and t_2 are free for x in F .

To implement the second observation, we add $\exists!z$ as an additional discharged assumption to ιE^1 :

$$\iota E^1 : \frac{\iota x[F, G] \quad \underbrace{F_z^x, G_z^x, \exists!z}_\Pi \quad C}{C} i$$

where z is not free in C nor any undischarged assumptions it depends on except F_z^x , G_z^x and $\exists!z$, and either z is the same as x or it is not free in F nor in G .

To find suitable modifications of the introduction rule for ι , we look at the derivation of $\iota x[F, G]$ from $\exists x(F \wedge \forall y(F_y^x \rightarrow x = y) \wedge G)$ in **I**. Had the application of the universal quantifier elimination rule been one of free logic, a further assumption $\exists!y$ would have been required, and had the existential quantifier elimination rule been one of free logic, it would have allowed the discharge of an assumption $\exists!x$. The latter lends itself as an additional premise of ιI , the former as an additional assumption discharged by that rule, which is again analogous to the existence assumptions required and discharged in applications of the rules for the quantifiers of free logic.

Generalising the second observation, we add $\exists!t$ as a further premise, and to implement the first observation we add $\exists!z$ as a further discharged assumption to ιI :

$$\iota I : \frac{\frac{F_t^x \quad G_t^x \quad \exists!t}{\iota x[F, G]} \quad \frac{\overbrace{F_z^x, \exists!z}^i}{z = t} \quad \Pi}{i}$$

where t is free for x in F and in G , and z is different from x , not free in t and does not occur free in any undischarged assumptions in Π except F_z^x and $\exists!z$.⁵

It is obvious that $\iota x[F, G]$ and $\exists x(F \wedge \forall y(F_y^x \rightarrow x = y) \wedge G)$ are interderivable in **INF** when ι is governed by the modified rules, but we give the deductions for convenience.

1. $\iota x[F, G] \vdash \exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)$

Let x and y be different variables, where y is not free in F nor in G :

$$\begin{array}{c} \iota x[F, G] \text{ --- } \frac{\overline{\exists!y}^2 \quad \overline{\exists!x}^3 \quad \overline{F_y^x}^1 \quad \overline{F}^3}{\iota E^2} \\ \frac{\frac{y = x}{(F_y^x \rightarrow y = x)}^1}{\forall y(F_y^x \rightarrow y = x)}^2 \quad \overline{F}^3 \\ \frac{F \wedge \forall y(F_y^x \rightarrow y = x) \quad \overline{G}^3}{F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G} \\ \frac{\iota x[F, G] \quad \frac{F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G}{\exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)}}{\exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)}^3 \iota E^1 \end{array}$$

⁵A more precise and general statement of the introduction rule for ι would result if we were to require Π to be a deduction of $(y = t)_z^y$ from $(F_y^x)_z^y$ and $\exists!z$, where y is different from x and not free in t , and either z is the same as y or z is not free in F_y^x nor in $y = t$.

2. $\exists x(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G) \vdash \iota x[F, G]$

Let \otimes be the formula $(F \wedge \forall y(F_y^x \rightarrow y = x) \wedge G)$, where y is different from x and not free in F or G :

$$\frac{\frac{\frac{\frac{\otimes}{F}^2}{\otimes}^2}{\exists!x}^2 \quad \frac{\frac{\frac{\frac{\otimes}{G}^2}{\otimes}^2}{\exists!y}^1}{\forall y(F_y^x \rightarrow y = x)}^2 \quad \frac{\frac{\otimes}{F_y^x}^1}{y = x}^1}{\iota I}^1 \quad \frac{\otimes}{\iota x[F, G]}^2}{\exists x \otimes}^2 \quad \iota x[F, G]$$

Let \mathbf{INF}^ι denote the systems of intuitionist negative free logic augmented with the rules for ι given in this section.

In order to prove a normalisation theorem for \mathbf{INF}^ι , we first observe that $\perp E$ can be restricted to atomic conclusions in this system:

1. Instead of inferring $\forall xA$ from \perp , infer A_y^x , for some y not occurring in any assumption that \perp depends on, and apply $\forall I$, discharging vacuously.
2. Instead of inferring $\exists xA$ from \perp , infer A_t^x , for some t that is free for x in A , infer $\exists!t$, and apply $\exists I$.
3. Instead of inferring $\iota x[F, G]$ from \perp , infer $F_t^x, G_t^x, \exists!t$ and $z = t$, for some t that is free for x in F and in G and some z that is not free in any assumption that \perp depends on, and apply ιI , discharging vacuously.

Next, $= E$ can be restricted to atomic formulas in \mathbf{INF}^ι . Consider an application of this rule with premise $\iota x[F, G]_{t_1}^y$:

$$\frac{t_1 = t_2 \quad \iota x[F, G]_{t_1}^y}{\iota x[F, G]_{t_2}^y}$$

where t_1 and t_2 are free for y in $\iota x[F, G]$. The exclusion of vacuous applications of $= E$ means that y must be different from x , and so $\iota x[F, G]_{t_1}^y$ is $\iota x[F_{t_1}^y, G_{t_1}^y]$. Let v and z be different variables not occurring in F, G, t_1, t_2 . The induction step applying $= E$ to subformulas of $\iota x[F, G]_{t_1}^y$ is the following:

$$\frac{\frac{\mathbf{A} \quad \mathbf{B} \quad \mathbf{C} \quad \mathbf{D}}{\iota x[F_{t_2}^y, G_{t_2}^y]} \quad \iota x[F_{t_1}^y, G_{t_1}^y]}{\iota x[F_{t_2}^y, G_{t_2}^y]} \quad \iota E^1,$$

where

$$\mathbf{A} = \frac{t_1 = t_2 \quad \overline{(F_{t_1}^y)^x}^2}{(F_{t_2}^y)^x} \quad \mathbf{B} = \frac{t_1 = t_2 \quad \overline{(G_{t_1}^y)^x}^2}{(G_{t_2}^y)^x} \quad \mathbf{C} = \overline{\exists!z}^2$$

$$\mathbf{D} = \frac{\iota x[F_{t_1}^y, G_{t_1}^y] \quad \overline{\exists!v}^1 \quad \overline{\exists!z}^2 \quad \overline{(F_{t_1}^y)^x}^1 \quad \overline{(F_{t_1}^y)^x}^2}{v = z} \quad \iota E^2.$$

As for applications of $\forall I$ and $\exists E$, we can assume that every application of ιI and ιE^1 has its own free variable, i.e. the variable z of an application of ιI or ιE^2 occurs only in the premises discharged by the rule and formulas derived from the discharged premises, and nowhere else in the deduction.

I will now give the reduction procedures for maximal formulas of the form $\iota x[F, G]$ and the permutative reduction procedures for maximal segments consisting of a formula of that form.

There are two cases of reduction procedures for maximal formulas of the form $\iota x[F, G]$ to be considered. First, the conclusion of ιI is the major premise of ιE^1 :

$$\frac{\frac{\frac{\Sigma_1 \quad \Sigma_2 \quad \Sigma_3}{F_t^x \quad G_t^x \quad \exists!t} \quad \frac{\overline{F_z^x}^i, \overline{\exists!z}^i}{z = t} \quad \Pi}{\iota x[F, G]} \quad \frac{\frac{\overline{F_v^x}^j, \overline{G_v^x}^j, \overline{\exists!v}^j}{\Xi} \quad C}{C} \quad j}{C} \quad i$$

Transform such steps in a deduction into the following, where Ξ_t^v is the deduction resulting from Ξ by replacing the variable v everywhere with the term t :

$$\frac{\frac{\Sigma_1 \quad \Sigma_2 \quad \Sigma_3}{F_t^x, G_t^x, \exists!t}}{\Xi_t^v} \quad C$$

First example. The major premise of ιE^1 is concluded by $\exists E$:

$$\frac{\frac{\frac{\overline{A_y^v}^i}{\Sigma} \quad \frac{\overline{F_z^x}^j, \overline{G_z^x}^j, \overline{\exists!z}^j}{\Pi}}{\exists v A \quad \iota x[F, G]}^i \quad C_j}{\iota x[F, G]}^i \quad C}{C}^j$$

Replace such steps in a deduction by:

$$\frac{\frac{\frac{\overline{A_y^v}^i}{\Sigma} \quad \frac{\overline{F_z^x}^j, \overline{G_z^x}^j, \overline{\exists!z}^j}{\Pi}}{\iota x[F, G]}^i \quad C_j}{\exists v A \quad C}^i \quad C}{C}^i$$

Second example. The major premise of ιE^2 is the conclusion of $\exists E$:

$$\frac{\frac{\frac{\overline{A_y^v}^i}{\Sigma} \quad \frac{\overline{\exists v A} \quad \iota x[F, G]}^i}{\iota x[F, G]}^i \quad \exists!t_1 \quad \exists!t_2 \quad F_{t_1}^x \quad F_{t_2}^x}{t_1 = t_2}}{t_1 = t_2}$$

Replace such steps in a deduction by:

$$\frac{\frac{\frac{\overline{A_y^v}^i}{\Sigma} \quad \frac{\overline{\exists v A} \quad \iota x[F, G]}^i \quad \exists!t_1 \quad \exists!t_2 \quad F_{t_1}^x \quad F_{t_2}^x}{t_1 = t_2}}^i \quad C}{t_1 = t_2}^i$$

The remaining cases are similar.

I am not counting ιE^2 as an introduction rule for $=$. There is no general way of removing formulas $t_1 = t_2$ concluded by ιE^2 and eliminated by $= E$, as the following illustrates:

$$\frac{\iota x[F, G] \quad \exists!t_1 \quad \exists!t_2 \quad F_{t_1}^x \quad F_{t_2}^x}{\frac{t_1 = t_2}{A_{t_2}^x} \quad A_{t_1}^x}$$

Thus there are no further maximal formulas to be considered in \mathbf{INF}^u . After the theorem, I will give an alternative second elimination rule for ι that avoids this problem.

We have the following:

THEOREM 1. *For any deduction Π of A from Γ in \mathbf{INF}^u there is a deduction of the same conclusion from some of the formulas in Γ that is in normal form.*

PROOF: By induction over the rank of proofs. The length of a segment is the number of formulas it consists of and its degree the number of logical constants in that formula. Let a maximal formula be a maximal segment of length 1. The *rank* of a deduction is the pair $\langle d, l \rangle$, where d is the highest degree of a maximal segment or 0 if there is none, and l is the sum of the lengths of maximal segments of highest degree. $\langle d, l \rangle < \langle d', l' \rangle$ iff either (i) $d < d'$ or (ii) $d = d'$ and $l < l'$. Applying the reduction procedures to a suitably chosen maximal segment of highest degree and longest length reduces the rank of a deduction. \square

We can reformulate the second elimination rule for ι to incorporate an application of Leibniz' Law instead of concluding with an identity:

$$\iota E^{2A} : \frac{\iota x[F, G] \quad \exists!t_1 \quad \exists!t_2 \quad F_{t_1}^x \quad F_{t_2}^x \quad A_{t_1}^x}{A_{t_2}^x}$$

A can be restricted to atomic formulas, an induction over the complexity of formulas showing that the general version with A a formula of any degree is admissible. Call the system resulting from \mathbf{INF}^u by replacing ιE^2 with ιE^{2A} $\mathbf{INF}^{u'}$.

ιE^2 and ιE^{2A} are interderivable in virtue of the rules for identity:

1. To derive ιE^{2A} , given premises $\iota x[F, G], \exists!t_1, \exists!t_2, F_{t_1}^x$ and $F_{t_2}^x$, derive $t_1 = t_2$ by ιE^2 and apply $= E$ to it and the premise $A_{t_1}^x$ to derive $A_{t_2}^x$.
2. To derive ιE^2 , let A be $t_1 = x$, so that $A_{t_1}^x$ is $t_1 = t_1$: derive it from $\exists!t_1$ by $= I^n$, apply ιE^{2A} to derive $A_{t_2}^x$, i.e. $t_1 = t_2$.

Thus \mathbf{INF}^u and $\mathbf{INF}^{u'}$ are equivalent.

In $\mathbf{INF}^{\prime\prime}$, steps in a deduction that conclude $t_1 = t_2$ by ιE^{2A} (with $t_1 = t_1$ as $A_{t_1}^x$) and using it as the identity in Leibniz' Law are redundant: ιE^{2A} can instead be applied with the premise and conclusion of Leibniz' Law. Such identities can therefore be removed from deductions, and we are now at liberty to count them amongst the maximal formulas.

If a maximal formula arises from introducing $\iota x[F, G]$ by ιI and eliminating it by ιE^{2A} , we have the following situation:

$$\begin{array}{c}
 \underbrace{F_z^x, \exists! z^{-i}} \\
 \Sigma_1 \quad \Sigma_2 \quad \Sigma_3 \quad \Pi \\
 \frac{F_{t_1}^x \quad G_{t_1}^x \quad \exists! t_1 \quad z = t_1}{\iota x[F, G]} \quad i \quad \Xi_1 \quad \Xi_2 \quad \Xi_3 \quad \Xi_4 \quad \Xi_5 \\
 \frac{\quad \quad \quad \quad \quad \quad \quad \quad \exists! t_2 \quad \exists! t_3 \quad F_{t_2}^x \quad F_{t_3}^x \quad A_{t_2}^x}{A_{t_3}^x}
 \end{array}$$

We now have two options for removing the maximal formula. We can proceed as previously: conclude $t_2 = t_3$ by an application of Leibniz' Law to the conclusions $t_2 = t_1$ of $\Pi_{t_1}^z$ and $t_3 = t_1$ of $\Pi_{t_2}^z$, and then apply Leibniz' Law once more with $A_{t_2}^x$ as further premise and $A_{t_3}^x$ as conclusion. Alternatively, we can first conclude $A_{t_1}^x$ from the conclusion $t_2 = t_1$ of $\Pi_{t_1}^z$ and $A_{t_2}^x$, and then conclude $A_{t_3}^x$ from $A_{t_2}^x$ and the conclusion $t_3 = t_1$ of $\Pi_{t_2}^z$. Thus deductions in the system resulting by replacing ιE^2 by ιE^{2A} also normalise, and it has the additional advantage of avoiding identities concluded by ιE^2 and eliminated by Leibniz' Law.

Thus we have the following:

THEOREM 2. *For any deduction Π of A from Γ in $\mathbf{INF}^{\prime\prime}$ there is a deduction of the same conclusion from some formulas in Γ that is in normal form.*

Deductions in $\mathbf{INF}^{\prime\prime}$ have slightly neater proof-theoretic properties than those in \mathbf{INF}^{\prime} , as deductions in normal form in $\mathbf{INF}^{\prime\prime}$ do not contain redundant identities introduced by ιE^2 and eliminated by $= E$. Deductions in \mathbf{INF}^{\prime} are, however, slightly simpler if we are interested in establishing identities, and this will be the case if we are interested in comparing the present system with the standard treatment of ι as a term forming operator: axioms and rules for the latter invariably appeal to identity.

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